

# CS520 Theory of Programming Languages

(Theory Area)

Phd. Qualifying Exam (January 12, 2019) - OPEN BOOK

## Problem 1 (40 points)

Consider the syntax of the language  $E$  for the “palm calculator”.

$E$	$\rightarrow$	$n$	number
		$x$	variable
		$E + E$	addition
		$E - E$	subtraction
		$( E )$	parenthesis
		let $x = E$ in $E$	binding

We will implement this language  $E$  with  $\langle S, E, C \rangle$ -machine. The  $\langle S, E, C \rangle$ -machine is an abstract machine.  $S$  is a stack of values (ordered sequence).  $E$  is an environment (function  $: \text{Variable} \rightarrow \text{Value}$ ).  $C$  is a command sequence defined as following:

$C$	$\rightarrow$	add· $C$	/* add top two values and push the result back */
		sub· $C$	
		bind( $x$ )· $C$	/* bind $x$ in $E$ with the value of stack-top, and pop */
		unbind( $x$ )· $C$	/* delete the most-recent binding $x$ in $E$ */
		push( $x$ )· $C$	/* push the value of the variable $x$ */
		push( $n$ )· $C$	/* push the constant $n$ */
		$\varepsilon$	; empty command

- Define the denotational semantics of the language  $E$
- Define the transition semantics (abstract machine semantics) of the command  $C$ .  
If you need a new command, define its syntax and transition semantics freely.  
You may use your own function if it is “well-defined”.
- Define the compilation rules  $\triangleright$  from the program  $E$  to the command  $C$  sequence with inference rules

## Problem 2 (40 points)

Let  $\mathbf{P}$  be the vertical domain of the natural numbers,  $\{0, 1, 2, \dots, \infty\}$  where

$$0 \sqsubseteq 1 \sqsubseteq 2 \sqsubseteq 3 \sqsubseteq \dots \sqsubseteq \infty,$$

and  $\mathbf{P}'$  the two-element domain  $\{\perp, \top\}$  where  $\perp \sqsubseteq \top$ .

Then the monotone function  $f \ x \equiv \text{if } x = \infty \text{ then } \top \text{ else } \perp$  is not continuous.

- (a) Why is the function  $f$  “not continuous”? Explain.
- (b) Let a function  $f_k$  be  $f_k x \equiv \text{if } x \leq k \text{ then } \perp \text{ else } \top$ . Then, is the function  $f_k$  continuous? Explain why or why not?
- (c) Does a sequence  $f_0 f_1 f_2 \dots$  constitute a chain? Why or why not? What is the least element among  $f_k$  where  $k = 0, 1, \dots, \infty$ ?
- (  $f \sqsubseteq g$  iff  $\forall x \in \mathbf{P}. f x \sqsubseteq g x$  )

**Problem 3 (20 points)**

The set  $\mathcal{T}$  is defined to consist only of  $t$  that are inductively defined as follows:

$$t \rightarrow \cdot \mid / t , t /$$

Every  $t \in \mathcal{T}$  enjoys several invariant conditions on the numbers of commas and slashes in  $t$ . Prove by induction that the number of slashes in  $t$  is  $(2 * \text{the number of commas})$  for any element of  $\mathcal{T}$ . First, state your theorem formally (with some mathematical notation), then prove your theorem by induction.